

Compilers and Tools for Software Stack Optimisation

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Outline

Compilers for a Set-Top-Box

Compilers Potential

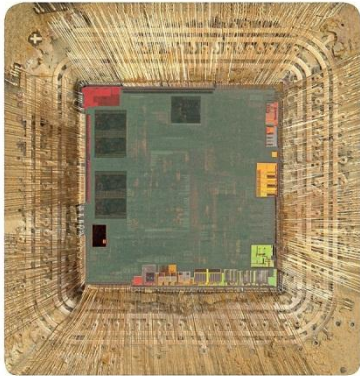
Auto Tuning Tools

Dynamic Program instrumentation

System calls interposition

COMPILERS FOR A SET-TOP-BOX

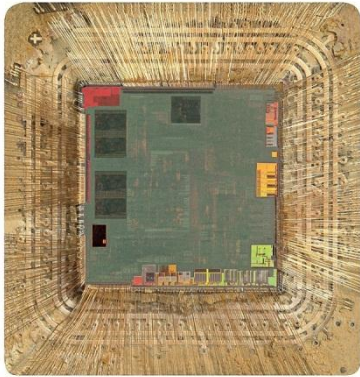
Set-Top-Box at Home



System on Chip (SoC)

Video Encode & Decode
Audio Encode & Decode
Connectivity and Storage
Drivers and Application Middlewares
Rich Content Applications

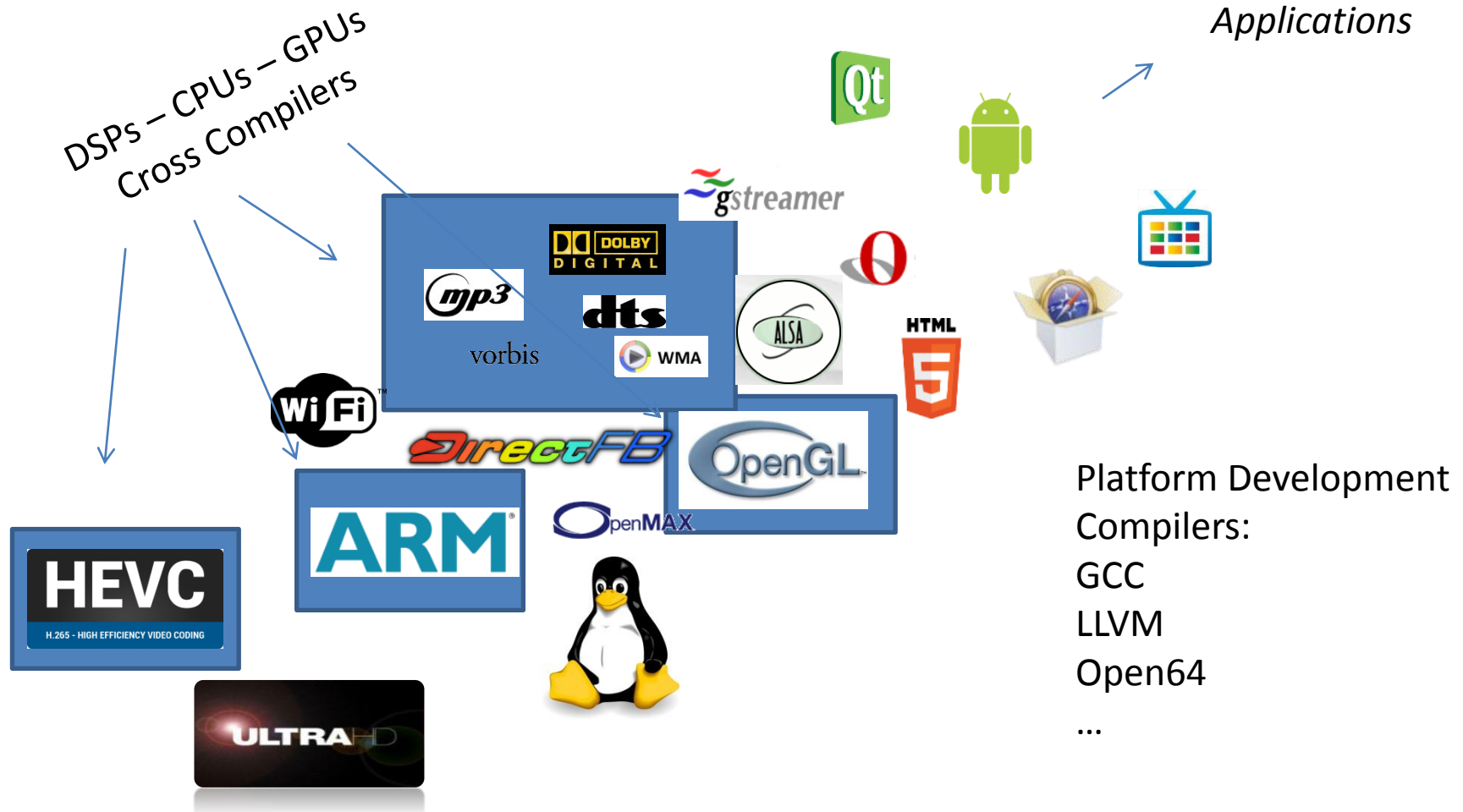
Set-Top-Box System on Chip



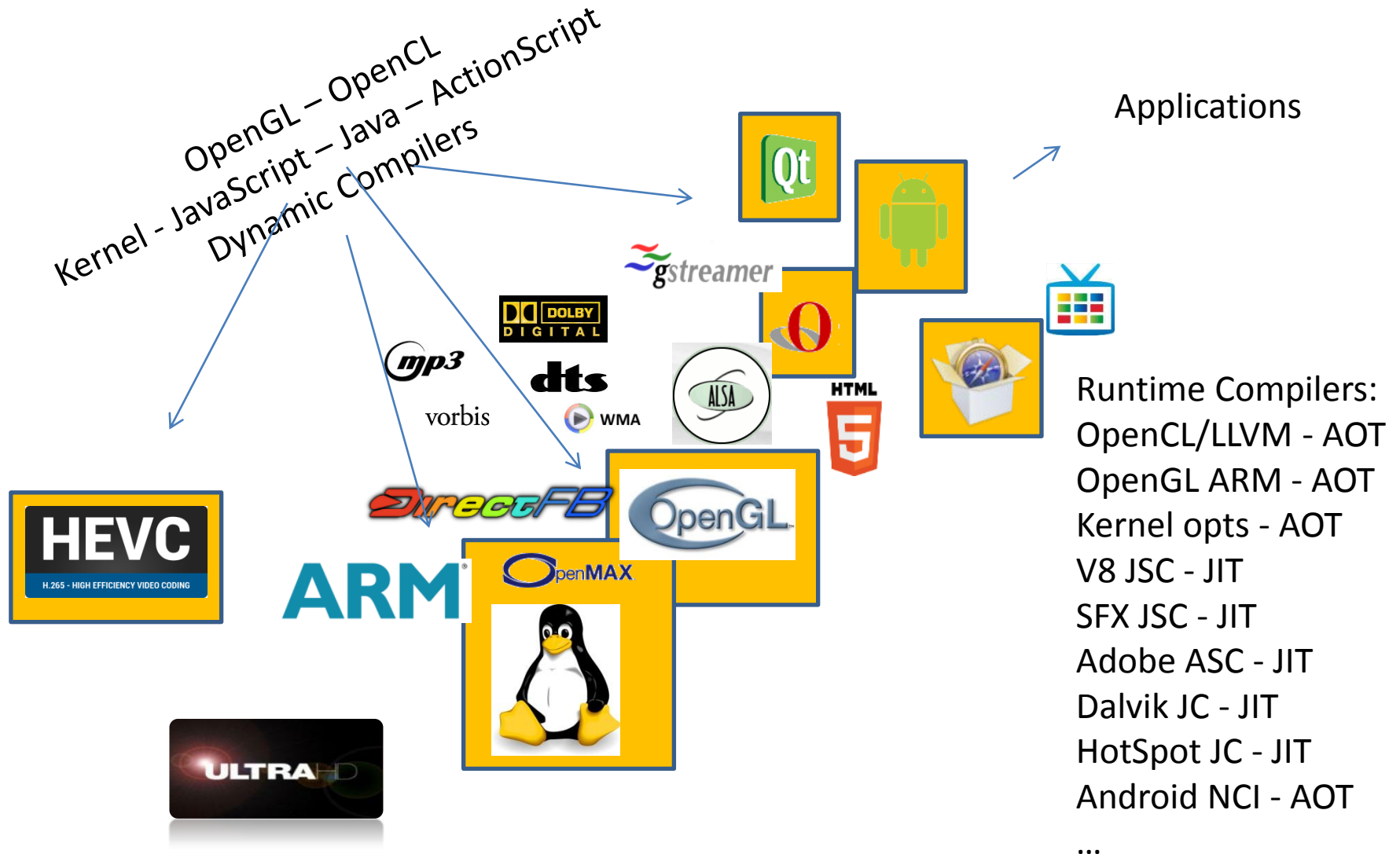
Applications



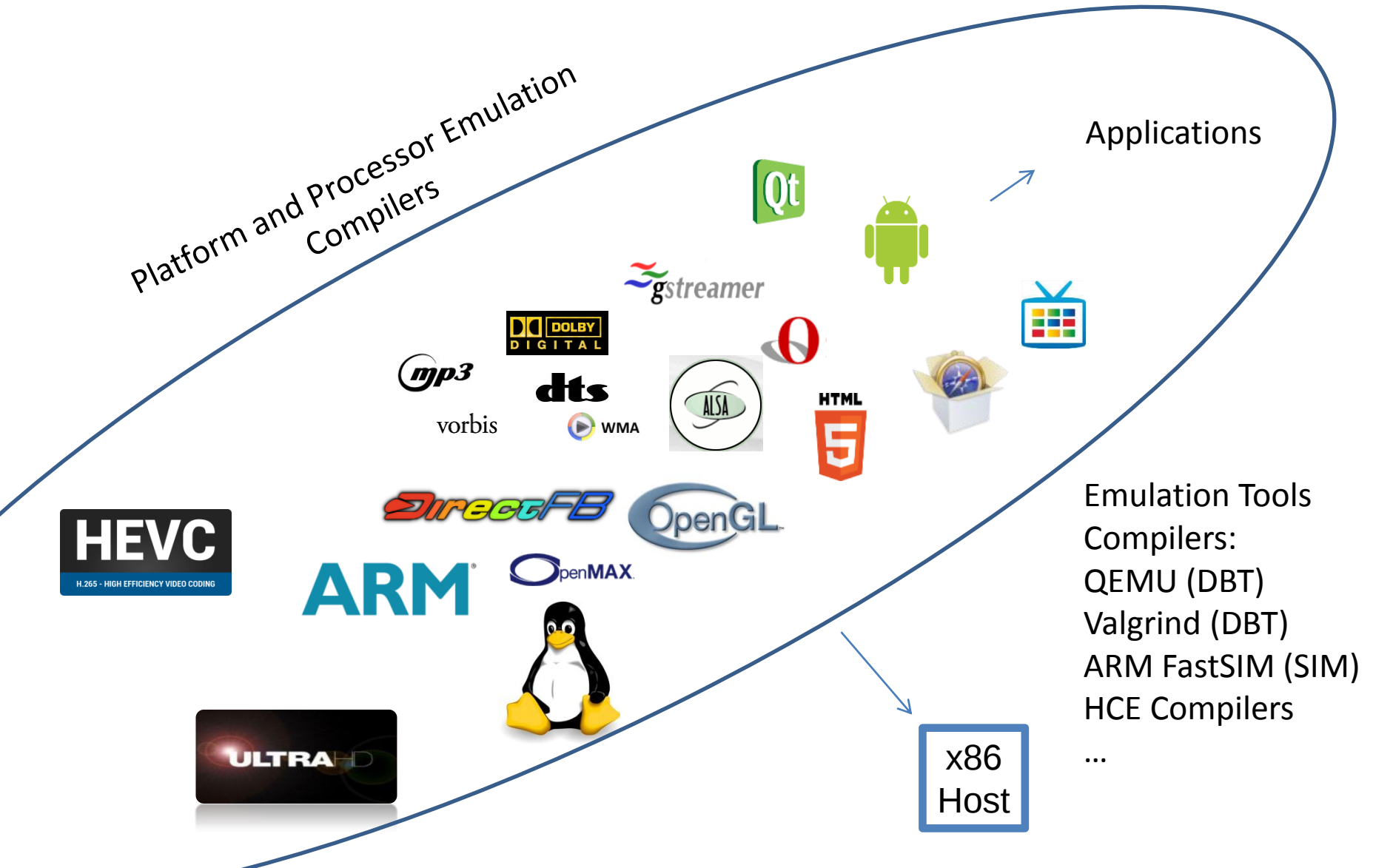
Cross Compilers



Embedded Compilers



Other Compilers



COMPILERS POTENTIAL

Compilers Potential and Cost

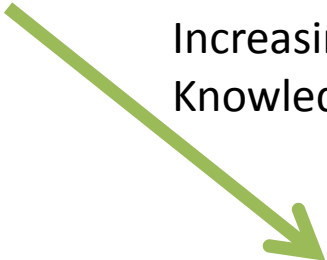
When?

Static

Ahead of Time

Just In Time

Increasing
Knowledge



Increasing
Runtime Cost



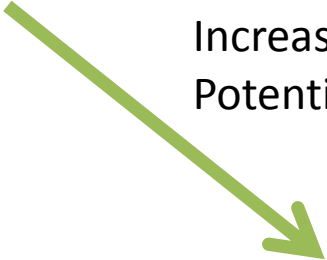
How Often?

One shot

Iterative

Continuous

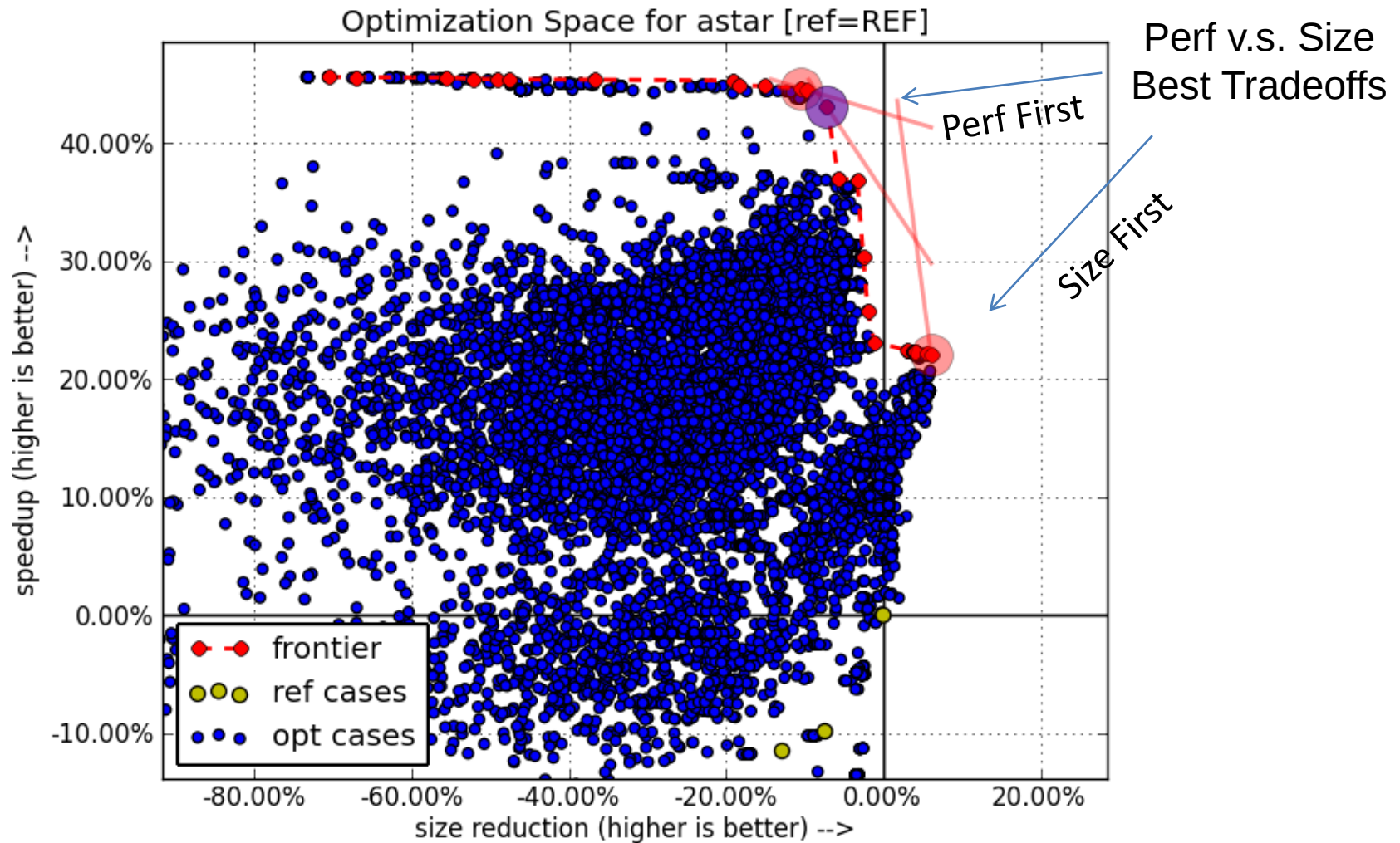
Increasing
Potential



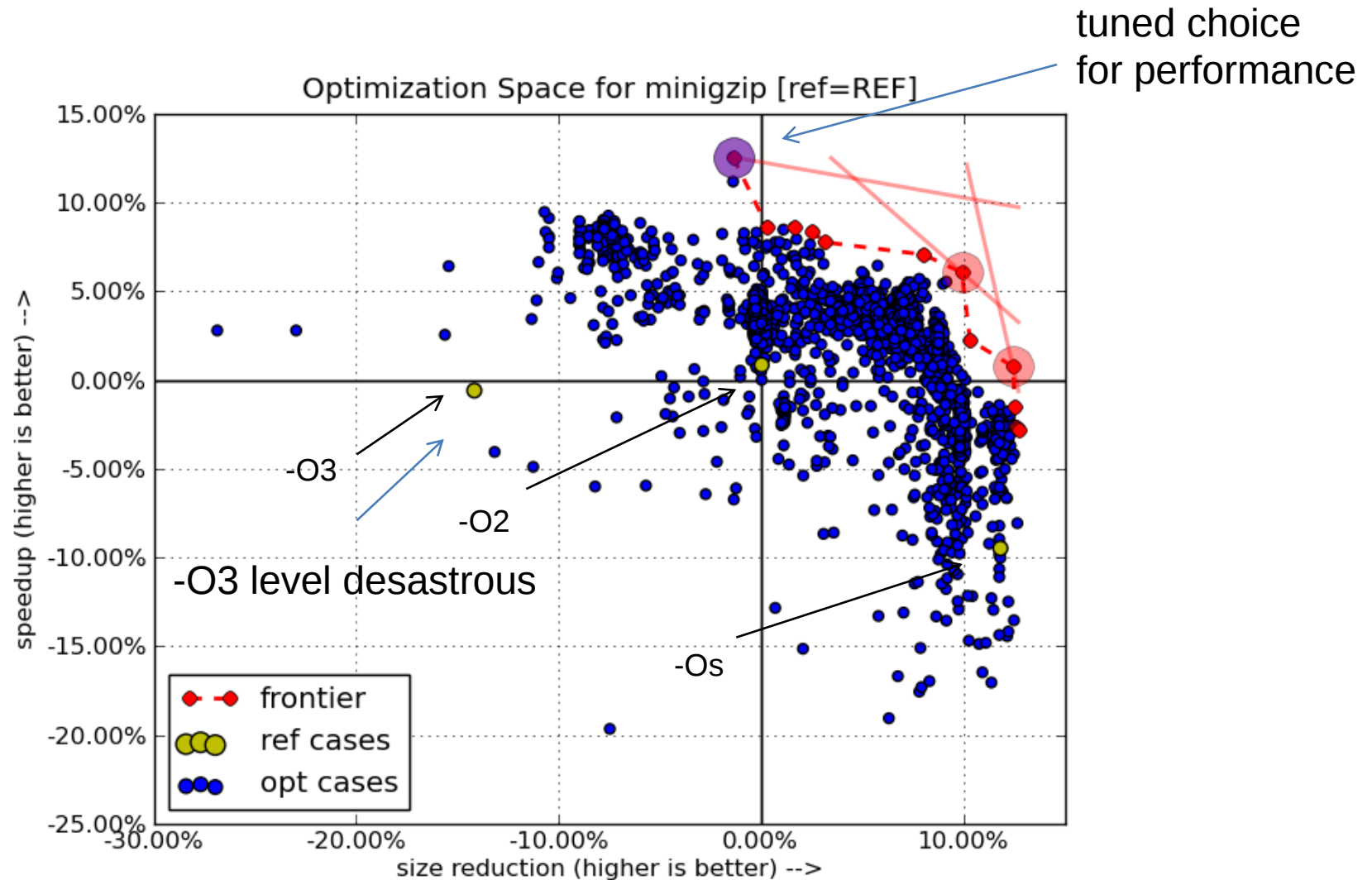
Increasing
Potential Cost



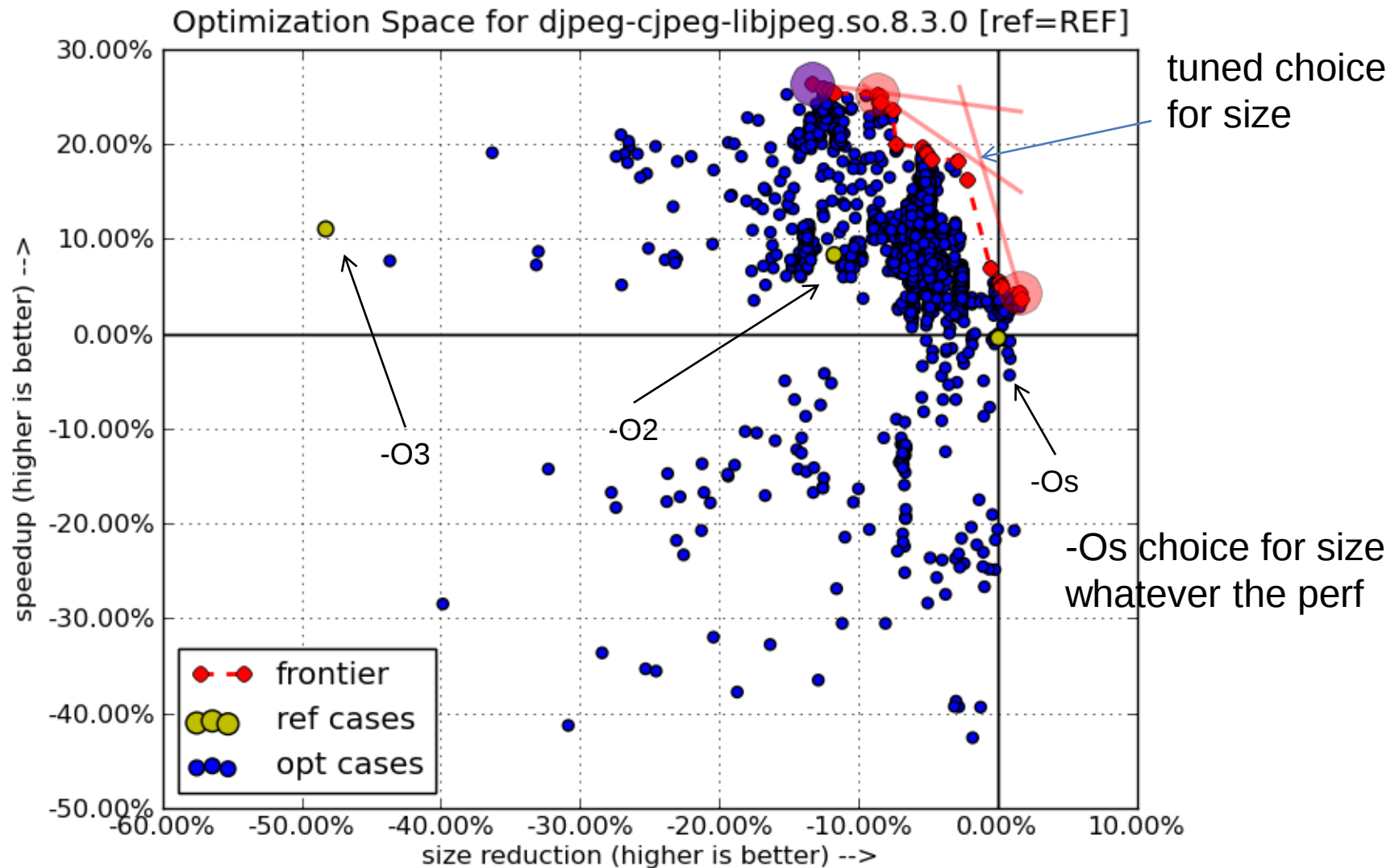
Optimizations Space Exploration



Defaults are Uneffective



Defaults Choices are Limited



Optimizing Applications

Choose the right domain specific language

Fallback to C/C++ 😊

Ensure Functionality and Unitary Testing at First

Develop Representative Benchmarks

Profile against the Benchmarks

Use your Domain Expertise

Use your Target Architecture Expertise 😞

Exploit the Compiler potential 😞

Iterate, but preserve maintainability

} Should you care?

AUTO TUNING TOOLS

Motivations

Automate the process of optimization

Build, run, observe, iterate, ...

Complement programmer's optimisation

Reproduce as changes are made

Explore more routes

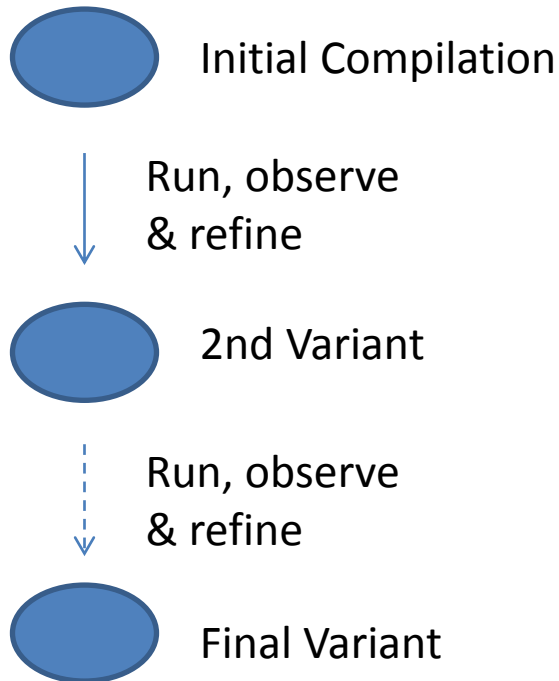
Exploit knowledge of compilers and architectures

Explore transparently

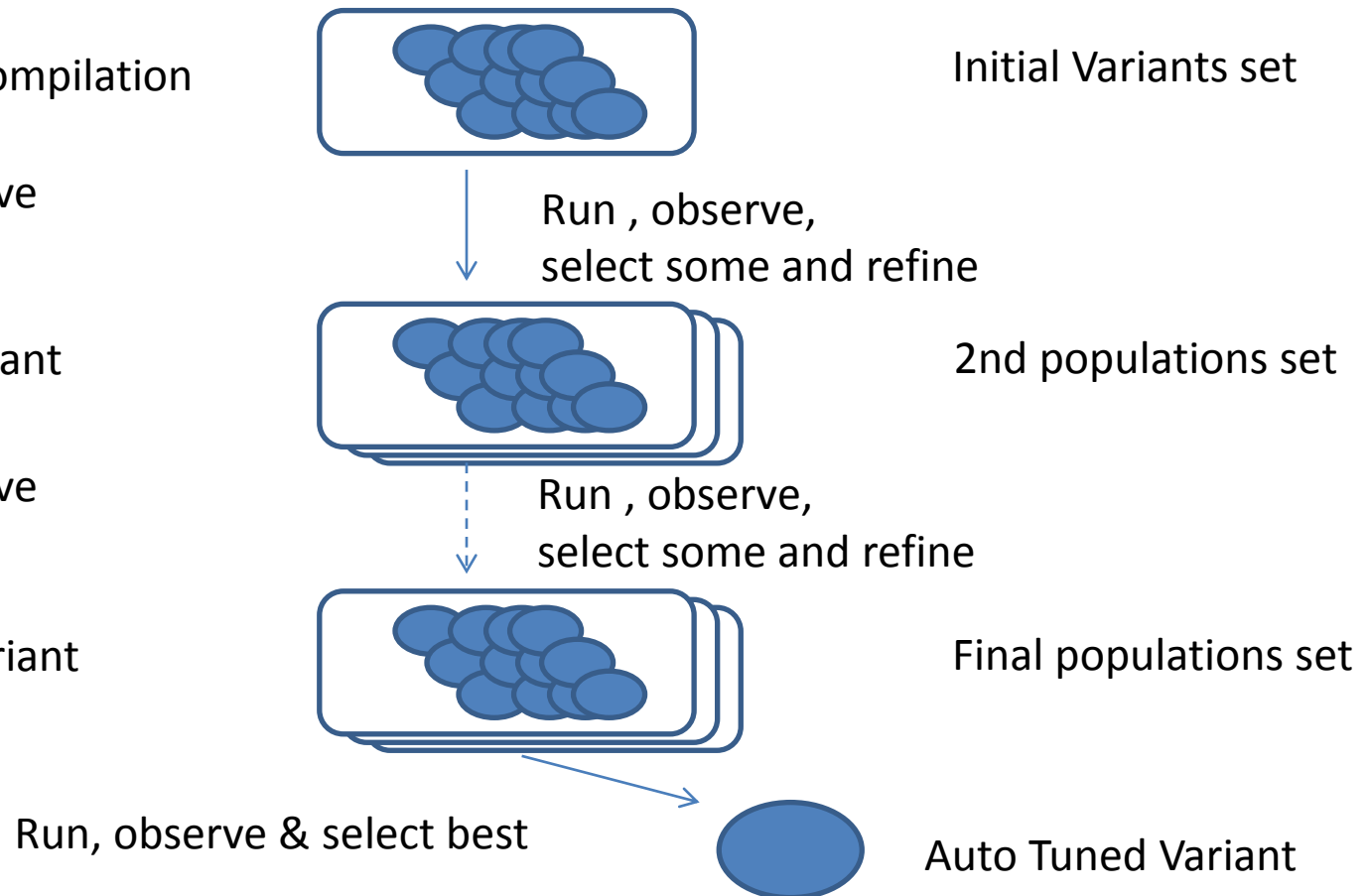
Inject optimisations without sources modification

Auto Tuning through Iterative Compilation

Iterative Compilation



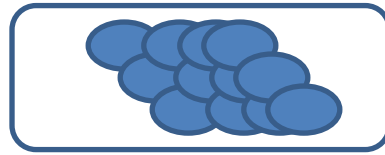
Multi variant Iterative Compilation



Cost of Auto Tuning

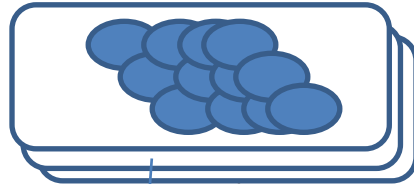
Populations sets (size P)

1st: base
(Ox X fdo X lto)



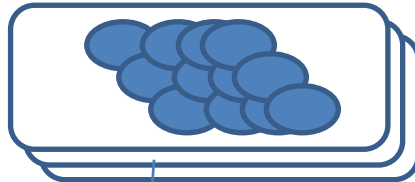
↓ Select 3

2nd: inlining
variants



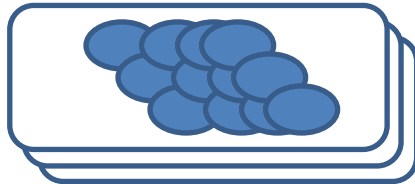
↓ Select 3

3rd: loop transfo
variants



↓ Select 3

4th: back-end
variants

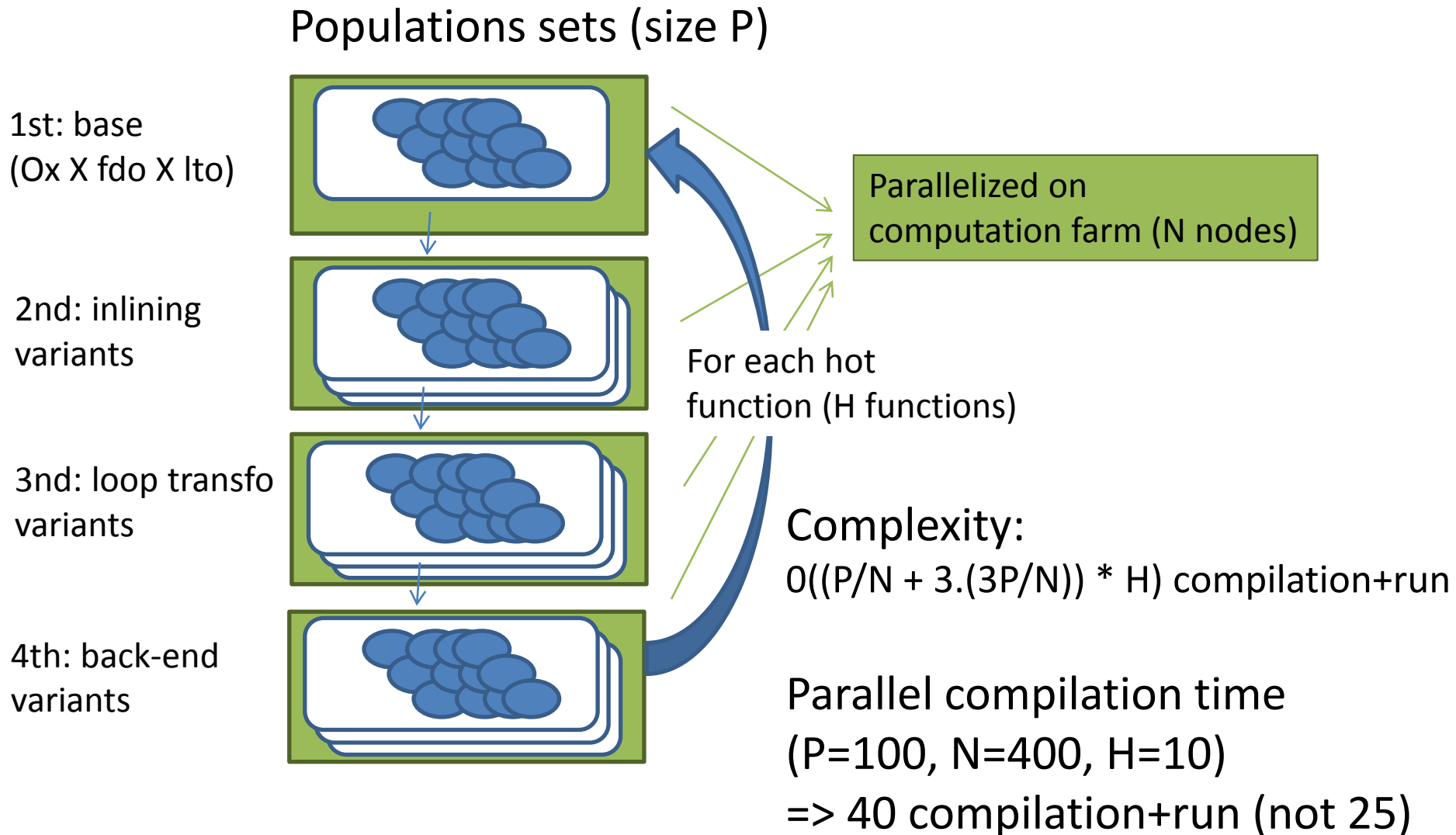


For each hot
function (H functions)

Complexity:
 $O(10.P.H)$ compilation+run

Typical total time cost
($P=100$, $H=10$):
 $\Rightarrow 10\,000$ compilation+run

Parallelizing Auto Tuning



Tools for Auto Tuning

Achieving Seamless Auto Tuning on Computation Farm

Compilation: audit build & produce variants *

Run: on chip or emulation **

Profiling: instrumentation ***

Observation: system perf/size ****

For instance

* with Proot + CARE

** with QEMU for instance

*** with QEMU + dynamic instrumentation

**** with bin/size and bin/perf tools for instance

Demo & Refs

Ref to auto tuning video:

<http://youtu.be/I5qpWGIajUs>

Auto tuning frameworks

TACT

ATOS (not yet disclosed)

Genetic exploration Tools

DYNAMIC PROGRAM INSTRUMENTATION

Motivations

Program performance analysis

cycles count, profiles, call graphs, ...

Profile driven compiler optimizations

l-cache placement, edge profiling, data dep. estimation, ...

Program debugging

call traces, syscall traces, memory checks, ...

Processor architecture analysis

instructions usage, cache behavior, ...

Dynamic Binary Translation

Program Loader

Disassembler

Compiler

Target Buffer Manager

Virtual Memory Manager

System calls Emulator

Debugger Bridge

Instrumentation Interface

Processes, Threads and Filesystem Monitor

Example with QEMU

QEMU, a versatile open source tool used for:

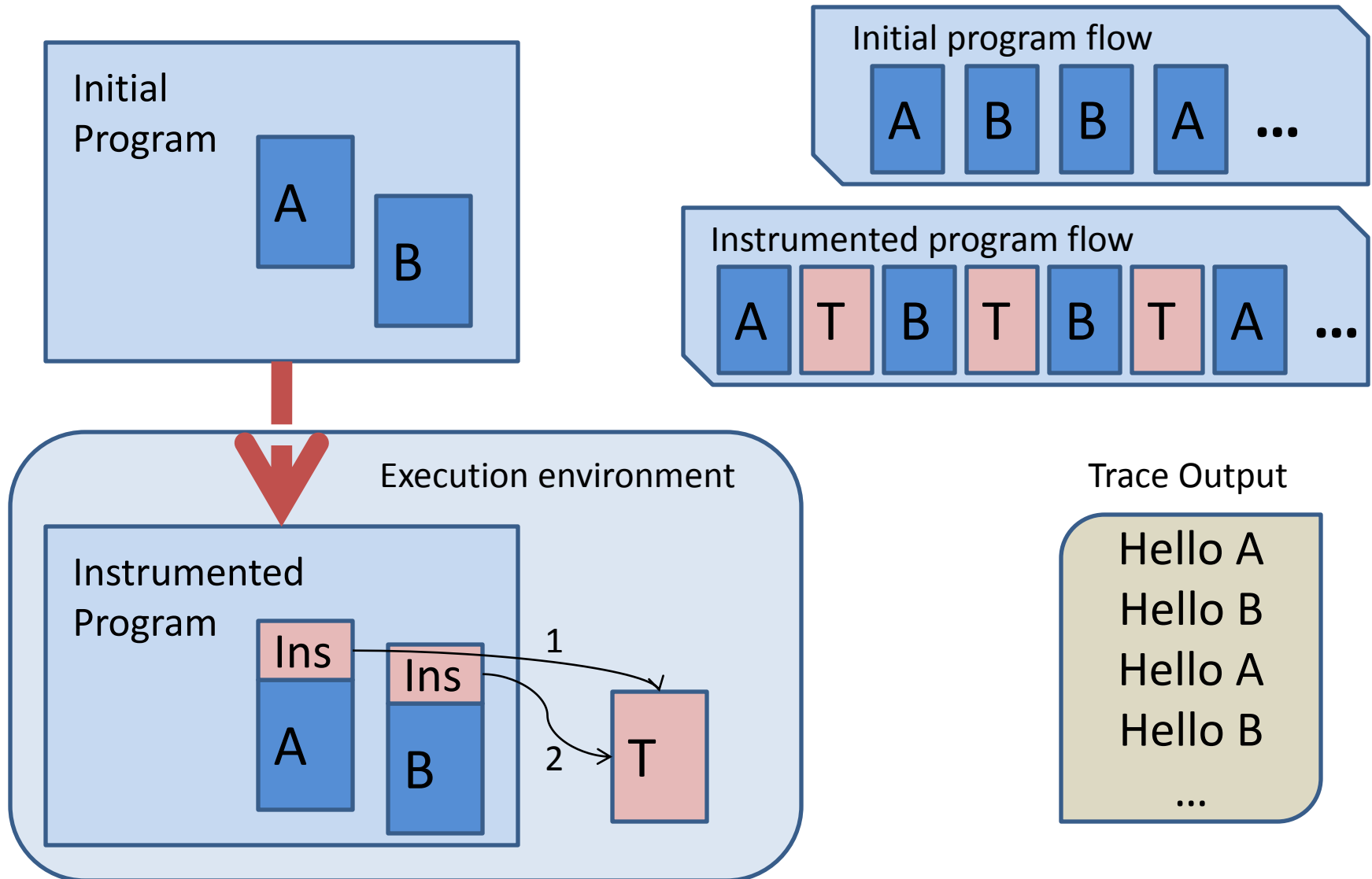
Platform emulation	(Google SDK,...)
Devices emulation	(VirtualBox,...)
Program emulation	(ScratchBox,...)

QEMU is a Dynamic Binary Translator (DBT)

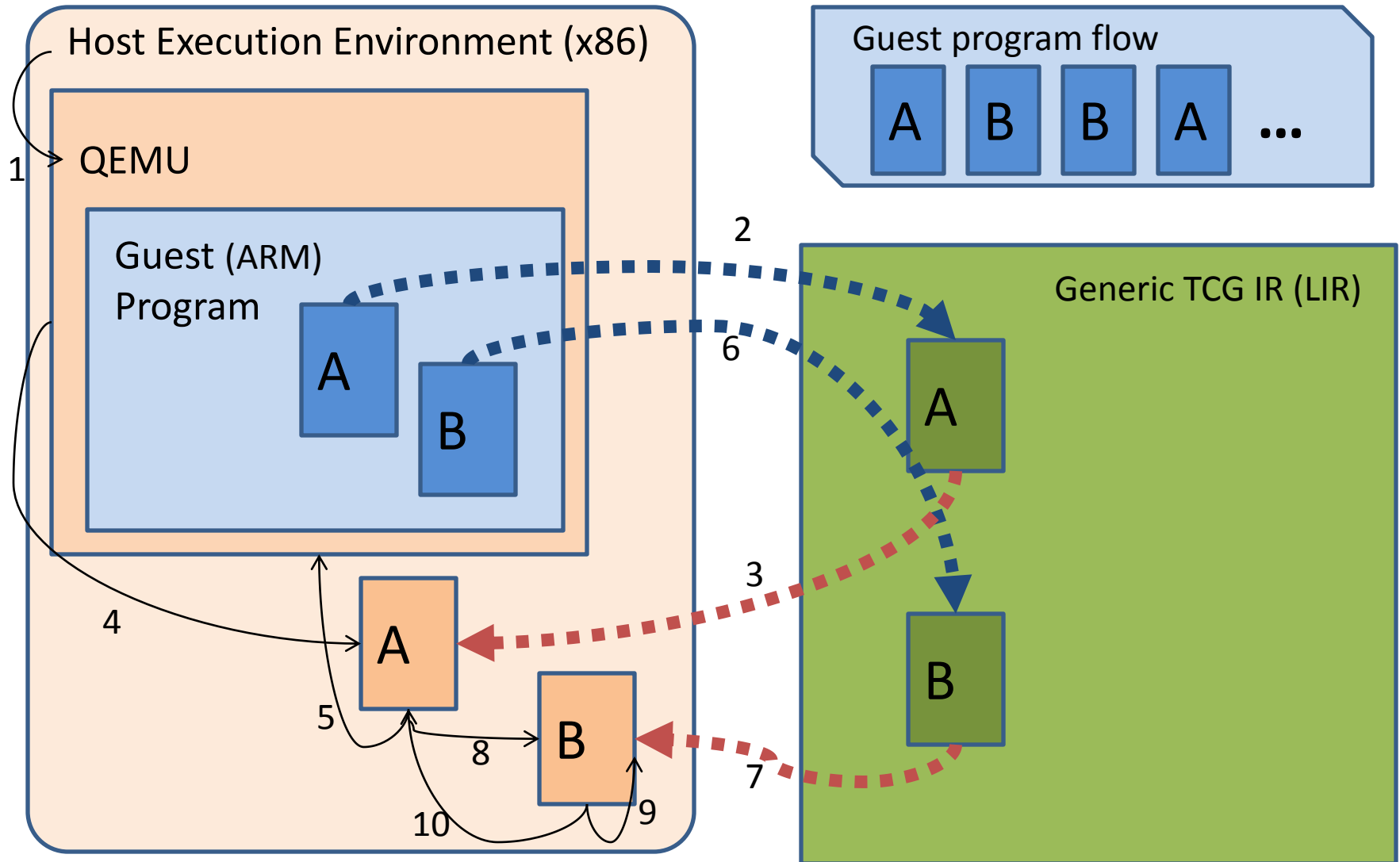
Hence it contains a compiler: TCG

(Tiny Code Generator)

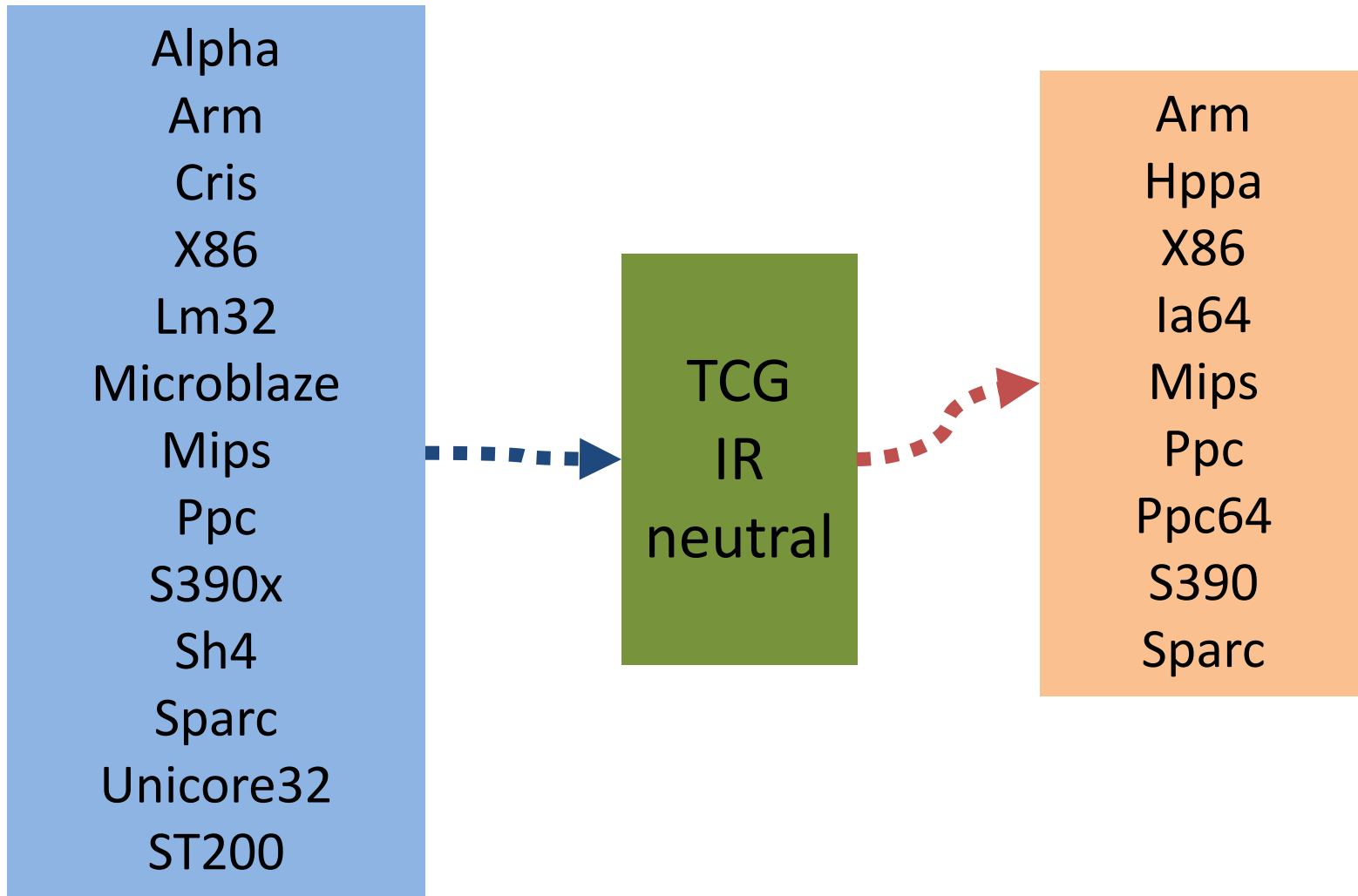
Program Instrumentation



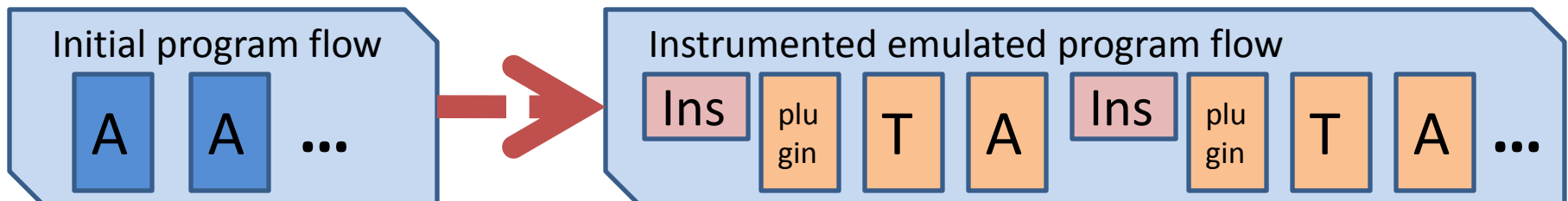
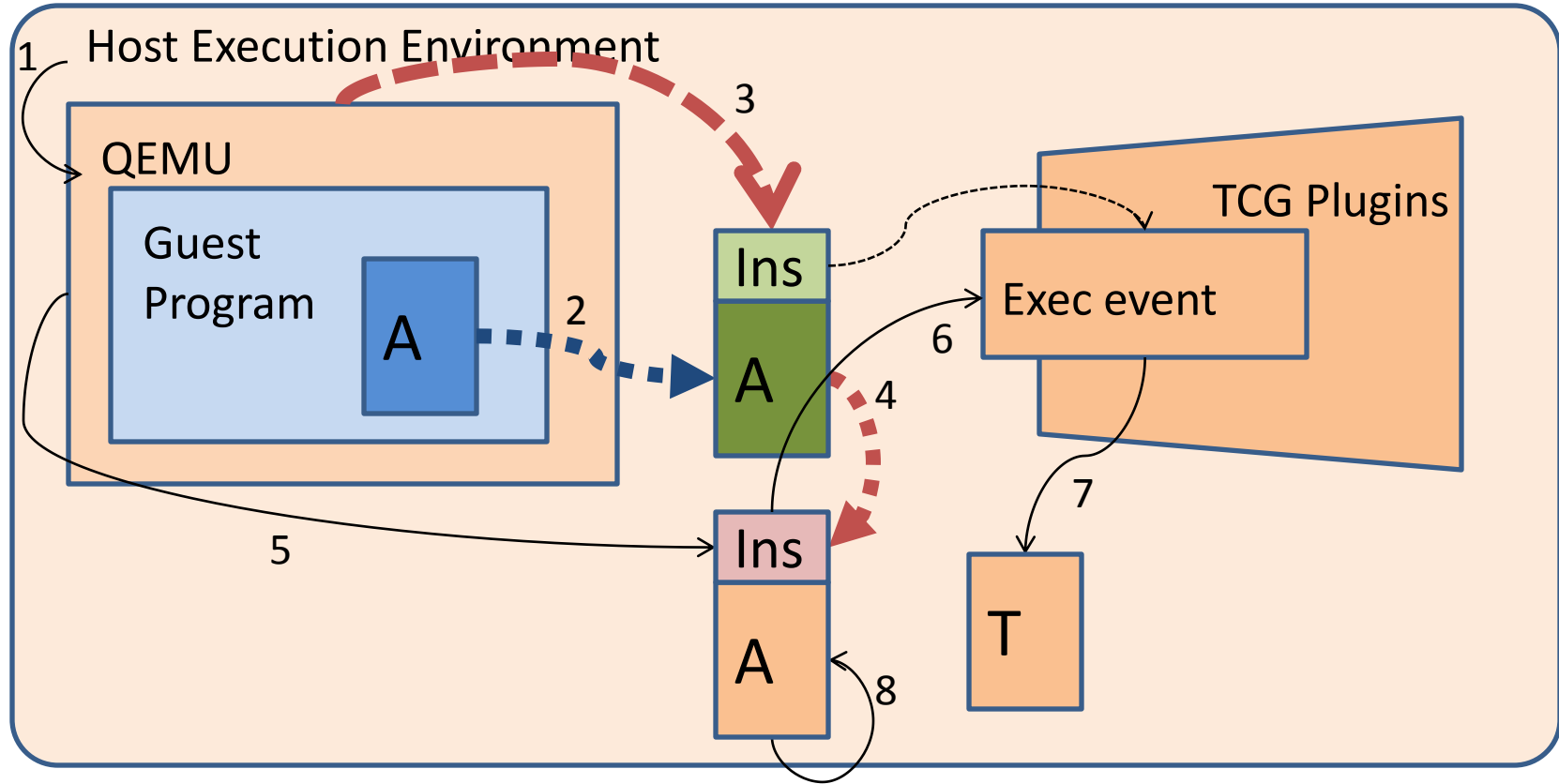
Binary Translation (i.e. QEMU ARM -> x86)



QEMU guests & hosts



Execution Time Instrumentation



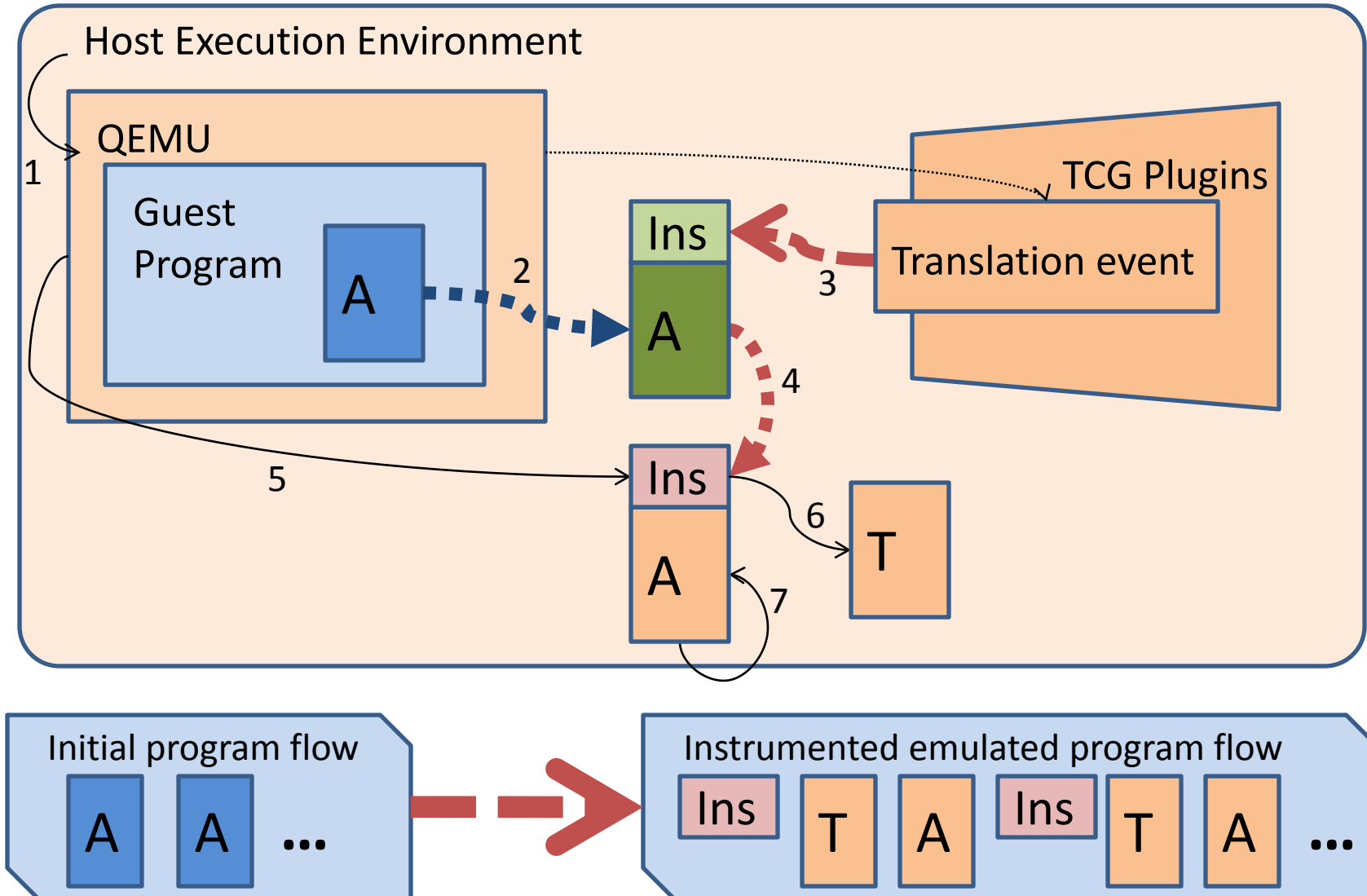
Instruction Count Plugin

```
/* A simple plugin for counting executed guest instructions.
 * Usage : cc shared -o icount.so icount.c
 * qemu-i386 -tcg-plugin ./icount.so the_program
 */
#include <stdint.h>
#include <stdio.h>
#include « qemu_tpi.h »

uint64_t total_icount;

/* Instruction count updated at each block execution. */
void qemu_tpi_block_execution_event(qemu_tpi_t *tpi) {
    total_icount += TPI_tb_icount(tpi);
}
void qemu_tpi_fini(qemu_tpi_t *tpi) {
    printf("Instructions: %"PRIu64"\n", total_icount);
}
```

Translation Time Instrumentation



Inlined Instruction Count Plugin

```
/* The translation event based code for the icount.c plugin. */

void qemu_tpi_block_translation_event(qemu_tpi_t *tpi) {
    /* Instruction count update inlined at each block translation.
     * Code is generated into the TCG IR buffer directly. */
    TPIv_ptr ptr = TPI_const_ptr(tpi, &total_icount);
    TPIv_i64 total = TPI_temp_new_i64(tpi);
    TPI_gen_ld_i64(tpi, total, ptr, 0);
    TPIv_i64 icount = TPI_const_i64(tpi, TPI_tb_icount(tpi));
    TPI_gen_add_i64(tpi, total, total, icount);
    TPI_temp_free_i64(tpi, icount);
    TPI_gen_st_i64(tpi, total, ptr, 0);
    TPI_temp_free_i64(tpi, total);
    TPI_temp_free_ptr(tpi, ptr);
}
```

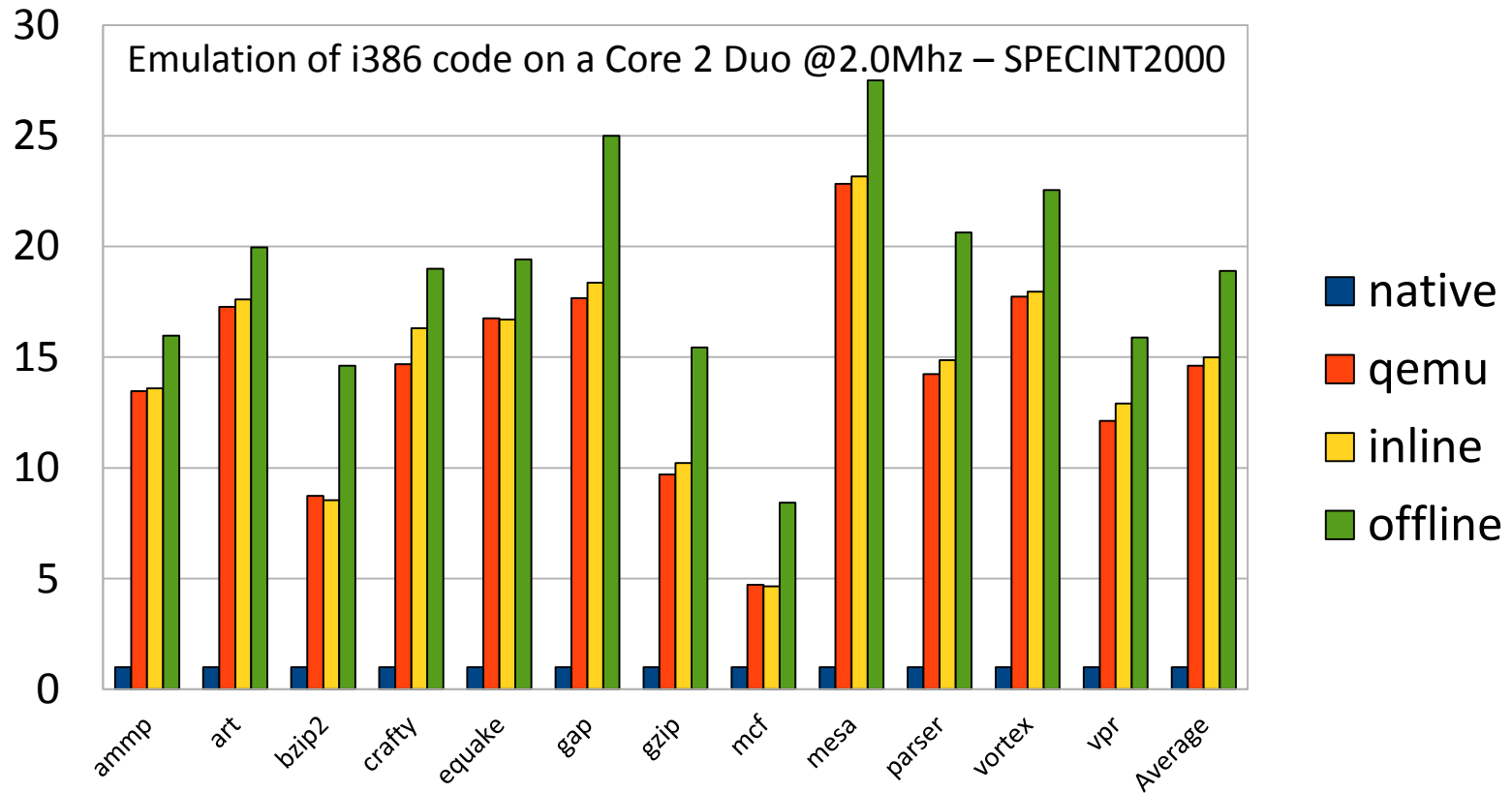
Profile Plugin Example

```
$ qemu-i386 -tcg-plugin profile.so sha1-386.exe
```

```
SHA1=15dd99a1991e0b3826fede3def fc1feba42278e6
```

Instrs	bytes	blocks	symbol
106497664	364229691	1792028	SHA1Transform
571297	1815640	133175	SHA1Update
961	3399	168	SHA1Final
23590	85199	6141	main
18	40	3	__libc_csu_init
16	43	5	.init
2093	12552	2087	.plt
55	146	10	.text
12	28	3	.fini
175102	504418	50840	unknown/libc.so.6

Instrumentation Overhead



Emulation cost is x15 (qemu)

Instrumentation overhead is 30% (offline) for a I-count plugin

Overhead reduced to 3% with translation time interface (inline)

Demo & Refs

Ref to live demo

Dynamic Binary Translators and
Instrumentation Tools

Valgrind

Pine

DynamoRIO

QEMU

SYSTEM CALLS INTERPOSITION

Motivation

Intercept communications between user programs and the kernel

Modify the behavior (side effects) of the user program without recompiling

Take full control (debugger)

Linux ptrace interface

Attach

fork(), ptrace(getpid(), ...) and exec(), or
ptrace(program_pid, ...)

Monitor

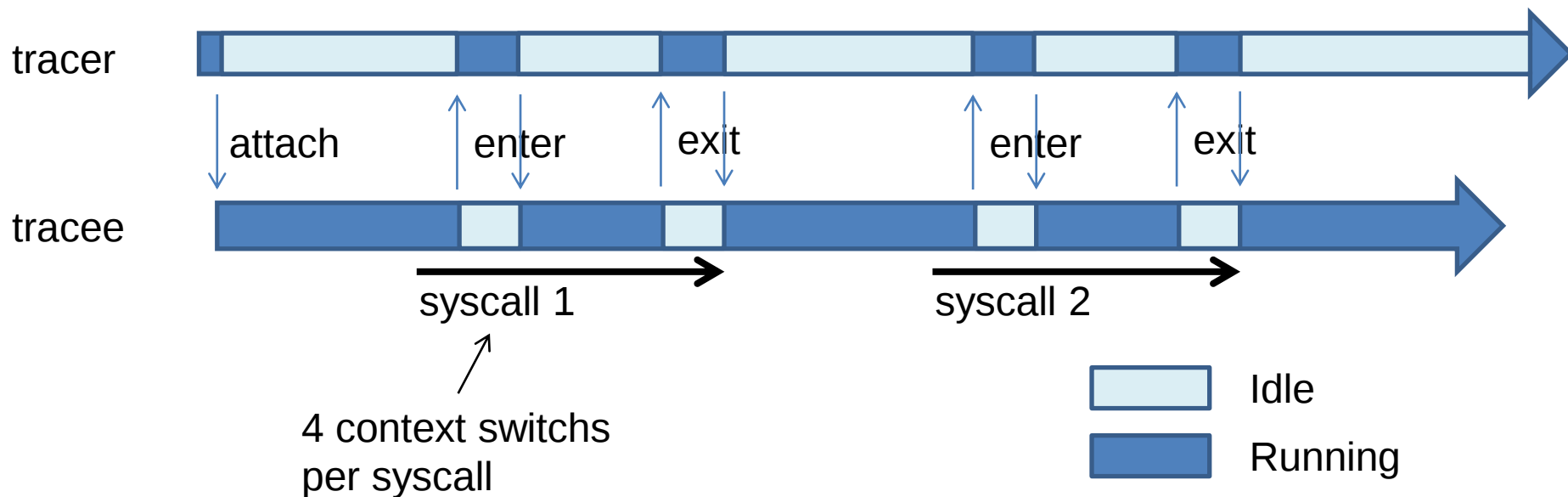
wait() and catch TRAP signal

Handle

get syscall() info and program state, and
modify syscall()/state and continue

Interposition and Context Switch

All syscall() from the tracee are notified
just before start
just before end



Interposition Example

Emulate a change of the `'/'` (root) location
equivalent to `chroot()` but in user-land
tracee behaves as in `'/'` but in your `/mydir`

Trace `open(/file_path)` (and all calls with paths)

Tracee: `open("/etc/passwd")`

Tracer: `"/etc/passwd" -> "/mydir/etc/passwd"`

Well, it's not that easy to fake passwords, why?

Also what about the loader, `"/mydir/lib/ld.so"` ?

Applications

Debugging

gdb, strace

Instrumentation

Valgrind

Cross filesystem execution

QEMU, Proot

Archiving

CARE, CDE

CARE

Comprehensive Archiving

For each path seen in a syscall

if (not_seen(path) && not_existing(path))

copy path to /archive/path

For Reproducible Execution

For each path seen in a syscall

transform path into /archive/path

Demo and Refs

See CARE demo video http://youtu.be/XIPzm66_3w0

CARE: <http://reproducible.io>

Proot: <http://proot.me>